

IV. Implementation

IV-A Method 1: Direct Implementation

以 STFT 為例

$$X(t, f) = \int_{-\infty}^{\infty} w(t - \tau) x(\tau) e^{-j2\pi f \tau} d\tau$$

Converting into the Discrete Form

$$t = n\Delta_t, \quad f = m\Delta_f, \quad \tau = p\Delta_t$$

$$X(n\Delta_t, m\Delta_f) = \sum_{p=-\infty}^{\infty} w((n-p)\Delta_t) x(p\Delta_t) e^{-j2\pi pm\Delta_t\Delta_f} \Delta_t$$

Suppose that $w(t) \cong 0$ for $|t| > B$, $B/\Delta_t = Q$

$$X(n\Delta_t, m\Delta_f) = \sum_{p=n-Q}^{n+Q} w((n-p)\Delta_t) x(p\Delta_t) e^{-j2\pi pm\Delta_t\Delta_f} \Delta_t$$

Problem : 對 scaled Gabor transform 而言 , $Q = ?$

- **Constraint for Δ_t** (The only constraint for the direct implementation method)

To avoid the aliasing effect,

$$\Delta_t < 1/2\Omega, \quad \Omega \text{ is the bandwidth of ?}$$

There is no constraint for Δ_f when using the direct implementation method.

Four Implementation Methods

(1) Direct implementation

Complexity:

假設 t -axis 有 T 個 sampling points, f -axis 有 F 個 sampling points

(2) FFT-based method

Complexity:

(3) FFT-based method with recursive formula

Complexity:

(4) Chirp-Z transform method

Complexity:

(A) Direct Implementation

Advantage : simple, flexible

Disadvantage : higher complexity

(B) DFT-Based Method

Advantage : lower complexity

Disadvantage : with some constraints

(C) Recursive Method

Advantage :

Disadvantage :

(D) Chirp Z Transform

Advantage :

Disadvantage :

IV-B Method 2: FFT-Based Method

Constraints : (ii) $\Delta_t \Delta_f = 1/N$,

(iii) $N = 1/(\Delta_t \Delta_f) \geq 2Q + 1$: ($\Delta_t \Delta_f$ 是整數的倒數)

Standard form of the DFT $Y[m] = \sum_{n=0}^{N-1} y[n] e^{-j \frac{2\pi mn}{N}}$

$$X(n\Delta_t, m\Delta_f) = \sum_{p=n-Q}^{n+Q} w((n-p)\Delta_t) x(p\Delta_t) e^{-j 2\pi pm \Delta_t \Delta_f} \Delta_t$$

$$\Delta_t \Delta_f = 1/N$$

$$X(n\Delta_t, m\Delta_f) = \sum_{p=n-Q}^{n+Q} w((n-p)\Delta_t) x(p\Delta_t) e^{-j \frac{2\pi pm}{N}} \Delta_t$$

$$q = p - (n - Q) \rightarrow p = (n - Q) + q$$

$$X(n\Delta_t, m\Delta_f) = \Delta_t e^{j \frac{2\pi(Q-n)m}{N}} \sum_{q=0}^{2Q} w((Q-q)\Delta_t) x((q+n-Q)\Delta_t) e^{-j \frac{2\pi qm}{N}}$$

Note that the input of the N -point FFT should have N points (others are set to zero).

$$X(n\Delta_t, m\Delta_f) = \Delta_t e^{j\frac{2\pi(Q-n)m}{N}} \sum_{q=0}^{N-1} x_1(q) e^{-j\frac{2\pi qm}{N}}, \quad q = p - (n-Q) \rightarrow p = (n-Q) + q$$

where $x_1(q) = w((Q-q)\Delta_t)x((n-Q+q)\Delta_t)$
 $x_1(q) = 0$

for $0 \leq q \leq 2Q$,
 for $2Q < q < N$.

$$n-Q \leq n-Q+q \leq n+Q$$

$$Q \geq Q-q \geq -Q$$

$$k = q - Q$$

$$X(n\Delta_t, m\Delta_f) = \Delta_t e^{j\frac{2\pi(Q-n)m}{N}} DFT(x_1(q))$$

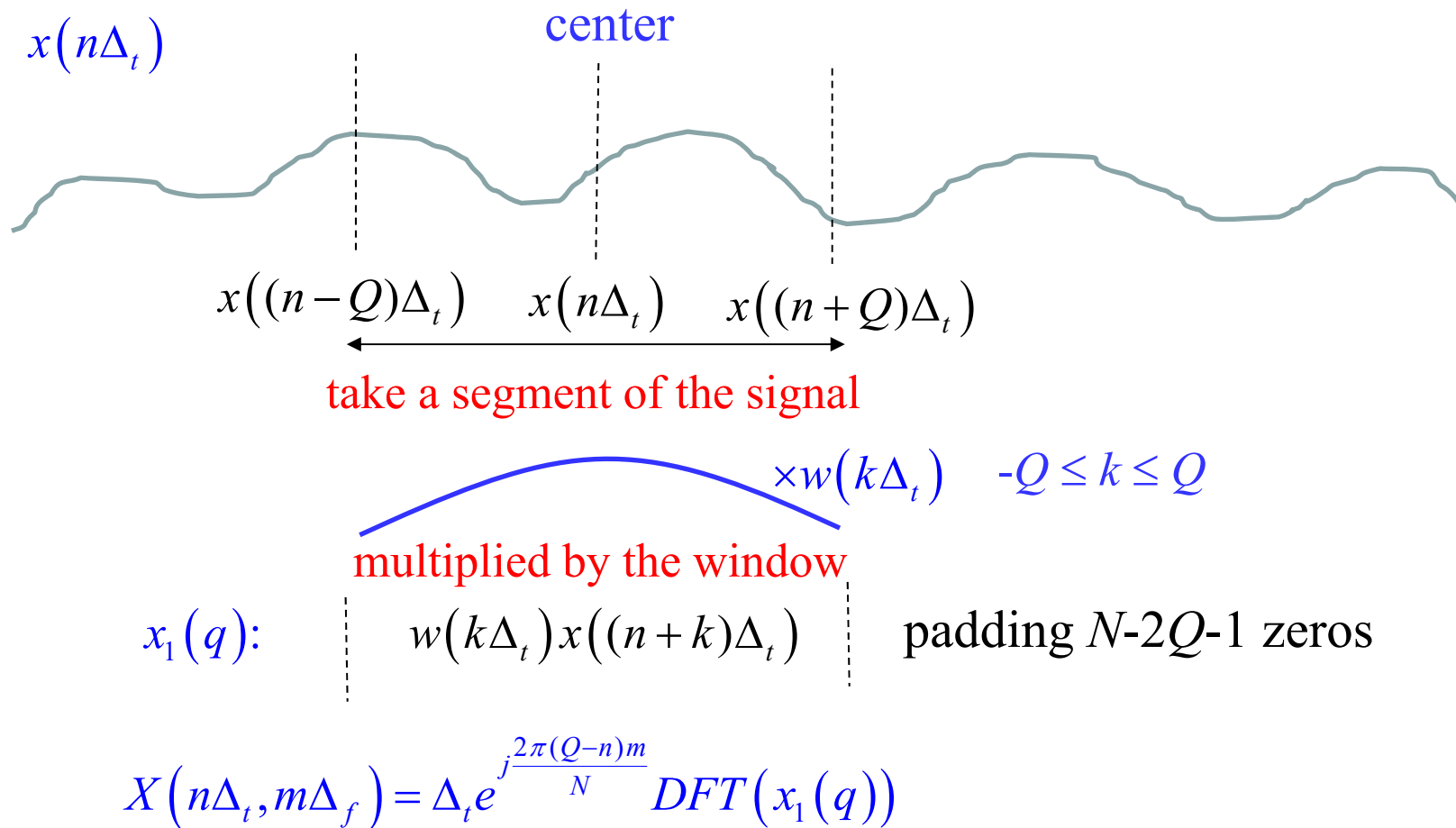
where $x_1(q) = w(k\Delta_t)x((n+k)\Delta_t)$ for $0 \leq q \leq 2Q$, $-Q \leq k \leq Q$ ($k = q - Q$)

$x_1(q) = 0$ for $2Q < q < N$.

(Suppose that $w(t) = w(-t)$)

$$X(n\Delta_t, m\Delta_f) = \Delta_t e^{j\frac{2\pi(Q-n)m}{N}} \text{DFT}(x_1(q))$$

where $x_1(q) = w(k\Delta_t)x((n+k)\Delta_t)$ for $-Q \leq k \leq Q$ ($k = q - Q$)
 $x_1(q) = 0$ for $2Q < q < N$. (Suppose that $w(t) = w(-t)$)



$$X(n\Delta_t, m\Delta_f) = \Delta_t e^{j\frac{2\pi(Q-n)m}{N}} \text{DFT}(x_1(q))$$

注意：

(1) 可以使用 Matlab 的 FFT 指令來計算 $\sum_{q=0}^{N-1} x_1(q) e^{-j\frac{2\pi qm}{N}} = X_1(m)$

(2) 對每一個固定的 n ，都要計算一次下方的式子

$$X(n\Delta_t, m\Delta_f) = \Delta_t e^{j\frac{2\pi(Q-n)m}{N}} \sum_{q=0}^{N-1} x_1(q) e^{-j\frac{2\pi qm}{N}} = \Delta_t e^{j\frac{2\pi(Q-n)m}{N}} X_1(m)$$

(fixed n)

(3) Complexity = ?

假設 $t = n_0\Delta_t, (n_0+1)\Delta_t, (n_0+2)\Delta_t, \dots, (n_0+T-1)\Delta_t$

$f = m_0\Delta_f, (m_0+1)\Delta_f, (m_0+2)\Delta_f, \dots, (m_0+F-1)\Delta_f$

Step 1: Calculate n_0, m_0, T, F, N, Q

Step 2: $n = n_0$

Step 3: Determine $x_1(q)$

Step 4: $X_1(m) = \text{FFT}[x_1(q)]$

Step 5: Convert $X_1(m)$ into $X(n\Delta_t, m\Delta_f)$

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$$m = f / \Delta_f$$

$$m_1 = \text{mod}(m, N) + 1$$

for Matlab

$$m_1 = m \% N$$

for Python

$$X(n\Delta_t, m\Delta_f) = X_1(?) \times ?$$

$$X_1[m] = \sum_{q=0}^{N-1} x_1(q) e^{-j \frac{2\pi qm}{N}}$$

$$X_1[m] = X_1[m + N]$$

Step 6: Set $n = n+1$ and return to Step 3 until $n = n_0+T-1$.

IV-C Method 3: Recursive Method

- A very fast way for implementing the rec-STFT

(n 和 $n-1$ 有 recursive 的關係)

$$X(n\Delta_t, m\Delta_f) = \sum_{p=n-Q}^{n+Q} x(p\Delta_t) e^{-j \frac{2\pi pm}{N} \Delta_t}$$

$$X((n-1)\Delta_t, m\Delta_f) =$$

- (1) Calculate $X(\min(n)\Delta_t, m\Delta_f)$ by the N -point FFT

$$X(n_0\Delta_t, m\Delta_f) = \Delta_t e^{j \frac{2\pi(Q-n_0)m}{N} N-1} \sum_{q=0}^{N-1} x_1(q) e^{-j \frac{2\pi qm}{N}}, \quad n_0 = \min(n),$$

$$x_1(q) = x((n-Q+q)\Delta_t) \quad \text{for } q \leq 2Q, \quad x_1(q) = 0 \quad \text{for } q > 2Q$$

- (2) Applying the recursive formula to calculate $X(n\Delta_t, m\Delta_f)$,

$$n = n_0 + 1 \sim \max(n)$$

$$X(n\Delta_t, m\Delta_f) = X((n-1)\Delta_t, m\Delta_f) - x((n-Q-1)\Delta_t) e^{-j2\pi(n-Q-1)m/N \Delta_t} \\ + x((n+Q)\Delta_t) e^{-j2\pi(n+Q)m/N \Delta_t}$$

T點 F點

IV-D Method 4: Chirp Z Transform

$$\exp(-j2\pi pm\Delta_t\Delta_f) = \exp(-j\pi p^2\Delta_t\Delta_f) \exp(j\pi(p-m)^2\Delta_t\Delta_f) \exp(-j\pi m^2\Delta_t\Delta_f)$$

For the STFT

$$X(n\Delta_t, m\Delta_f) = \sum_{p=n-Q}^{n+Q} w((n-p)\Delta_t) x(p\Delta_t) e^{-j2\pi pm\Delta_t\Delta_f} \Delta_t$$

$$X(n\Delta_t, m\Delta_f) = \Delta_t e^{-j\pi m^2\Delta_t\Delta_f} \sum_{p=n-Q}^{n+Q} w((n-p)\Delta_t) x(p\Delta_t) e^{-j\pi p^2\Delta_t\Delta_f} e^{j\pi(p-m)^2\Delta_t\Delta_f}$$

Step 1 multiplication

Step 2 convolution

Step 3 multiplication

$$\text{Step 1 } x_1[p] = w((n-p)\Delta_t)x(p\Delta_t)e^{-j\pi p^2\Delta_t\Delta_f} \quad n-Q \leq p \leq n+Q$$

$$\text{Step 2 } X_2[n, m] = \sum_{p=n-Q}^{n+Q} x_1[p]c[m-p] \quad c[m] = e^{j\pi m^2\Delta_t\Delta_f}$$

$$\text{Step 3 } X(n\Delta_t, m\Delta_f) = \Delta_t e^{-j\pi m^2\Delta_t\Delta_f} X_2[n, m]$$

Step 2 在計算上，需要用到 linear convolution 的技巧

Question: Step 2 要用多少點的 DFT?

- **Illustration for the Question on Page 124**

$$y[n] = \sum_k x[n-k]h[k]$$

- **Case 1**

When $\text{length}(x[n]) = N$, $\text{length}(h[n]) = K$, N and K are finite,

—————→ $\text{length}(y[n]) = N+K-1$,

Using the $(N+K-1)$ -point DFTs (學信號處理的人一定要知道的常識)

- **Case 2**

$x[n]$ has finite length but $h[n]$ has infinite length ????

$$y[n] = \sum_k x[n-k]h[k]$$

- Case 2

$x[n]$ has finite length but $h[n]$ has infinite length

$x[n]$ 的範圍為 $n \in [n_1, n_2]$ ，範圍大小為 $N = n_2 - n_1 + 1$

$h[n]$ 無限長

$y[n] = \sum_k x[n-k]h[k]$ $y[n]$ 每一點都有值 (範圍無限大)

但我們只想求出 $y[n]$ 的其中一段

希望算出的 $y[n]$ 的範圍為 $n \in [m_1, m_2]$ ，範圍大小為 $M = m_2 - m_1 + 1$

$h[n]$ 的範圍？

要用多少點的 FFT？

$$y[n] = x[n] * h[n] = \sum_{k=-\infty}^{\infty} x[n-k]h[k]$$

改寫成 $y[n] = x[n] * h[n] = \sum_{s=n_1}^{n_2} x[s]h[n-s]$

$$y[n] = x[n_1]h[n-n_1] + x[n_1+1]h[n-n_1-1] + x[n_1+2]h[n-n_1-2] \\ + \cdots + x[n_2]h[n-n_2]$$

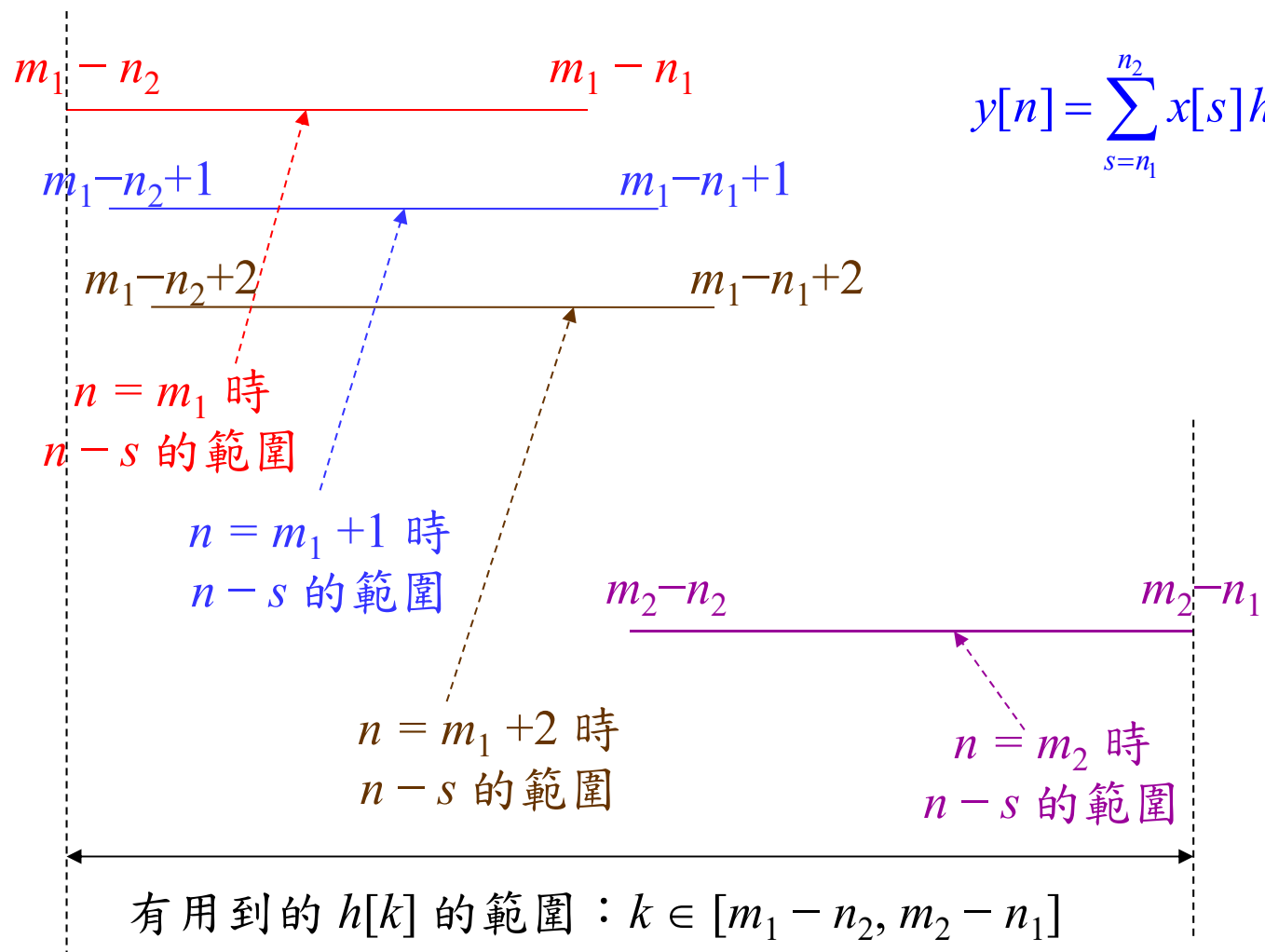
當 $n = m_1$

$$y[m_1] = x[n_1]h[m_1-n_1] + x[n_1+1]h[m_1-n_1-1] + x[n_1+2]h[m_1-n_1-2] \\ + \cdots + x[n_2]h[m_1-n_2]$$

當 $n = m_2$

$$y[m_2] = x[n_1]h[m_2-n_1] + x[n_1+1]h[m_2-n_1-1] + x[n_1+2]h[m_2-n_1-2] \\ + \cdots + x[n_2]h[m_2-n_2]$$

$$y[n] = \sum_{s=n_1}^{n_2} x[s]h[n-s]$$



所以，有用到的 $h[k]$ 的範圍是 $k \in [m_1 - n_2, m_2 - n_1]$

範圍大小為 $m_2 - n_1 - m_1 + n_2 + 1 = N + M - 1$

FFT implementation for Case 2

$$x_1[n] = x[n + n_1] \quad \text{for } n = 0, 1, 2, \dots, N-1$$

$$x_1[n] = 0 \quad \text{for } n = N, N+1, N+2, \dots, L-1 \quad L = N + M - 1$$

$$h_1[n] = h[n + m_1 - n_2] \quad \text{for } n = 0, 1, 2, \dots, L-1$$

$$y_1[n] = \text{IFFT}_L \left(\text{FFT}_L \{x_1[n]\} \text{FFT}_L \{h_1[n]\} \right)$$

$$y[n] = y_1[n - m_1 + N - 1] \quad \text{for } n = m_1, m_1+1, m_1+2, \dots, m_2$$

IV-E Unbalanced Sampling for STFT and WDF

將 pages 114 and 118 的方法作修正

$$X(t, f) = \int_{-\infty}^{\infty} w(t - \tau) x(\tau) e^{-j2\pi f \tau} d\tau$$

$$X(n\Delta_t, m\Delta_f) = \sum_{p=nS-Q}^{nS+Q} w((nS - p)\Delta_\tau) x(p\Delta_\tau) e^{-j2\pi pm\Delta_\tau\Delta_f} \Delta_\tau$$

where $t = n\Delta_t$, $f = m\Delta_f$, $\tau = p\Delta_\tau$, $B = Q\Delta_\tau$ (假設 $w(t) \cong 0$ for $|t| > B$),

$$S = \Delta_t / \Delta_\tau \quad \Delta_t \neq \Delta_\tau$$

註： Δ_τ (sampling interval for the **input** signal)

Δ_t (sampling interval for the **output t -axis**) can be different.

However, it is better that $S = \Delta_t / \Delta_\tau$ is an integer.

When (1) $\Delta_t \Delta_f = 1/N$, (2) $N = 1/(\Delta_t \Delta_f) > 2Q + 1$: ($\Delta_t \Delta_f$ 只要是整數的倒數即可)

(3) $\Delta_t < 1/2\Omega$, Ω is the bandwidth of $w(\tau - t)x(\tau)$

i.e., $|FT\{w(\tau - t)x(\tau)\}| = |X(t, f)| \approx 0$ when $|f| > \Omega$

$$X(n\Delta_t, m\Delta_f) = \sum_{p=nS-Q}^{nS+Q} w((nS-p)\Delta_t) x(p\Delta_t) e^{-j\frac{2\pi pm}{N} \Delta_t}$$

$$\text{令 } q = p - (nS - Q) \rightarrow p = (nS - Q) + q$$

$$X(n\Delta_t, m\Delta_f) = \Delta_t e^{j\frac{2\pi(Q-nS)m}{N}} \sum_{q=0}^{N-1} x_1(q) e^{-j\frac{2\pi qm}{N}}$$

$$x_1(q) = w((Q-q)\Delta_t) x((nS-Q+q)\Delta_t) \quad \text{for } 0 \leq q \leq 2Q,$$

$$x_1(q) = 0 \quad \text{for } 2Q < q < N.$$

If $w(t) = w(-t)$

$$x_1(q) = w(k\Delta_t) x((nS+k)\Delta_t) \quad \text{for } 0 \leq q \leq 2Q, \quad k = q - Q, \quad -Q \leq k \leq Q$$

$$x_1(q) = 0 \quad \text{for } 2Q < q < N.$$

$$\begin{aligned} \text{假設 } t &= c_0 \Delta_t, (c_0+1) \Delta_t, (c_0+2) \Delta_t, \dots, (c_0+C-1) \Delta_t \\ &= c_0 S \Delta_\tau, (c_0 S+S) \Delta_\tau, (c_0 S+2S) \Delta_\tau, \dots, [c_0 S+(C-1)S] \Delta_\tau \end{aligned}$$

$$f = m_0 \Delta_f, (m_0+1) \Delta_f, (m_0+2) \Delta_f, \dots, (m_0+F-1) \Delta_f$$

$$\tau = n_0 \Delta_\tau, (n_0+1) \Delta_\tau, (n_0+2) \Delta_\tau, \dots, (n_0+T-1) \Delta_\tau \quad S = \Delta_t / \Delta_\tau$$

Step 1: Calculate $c_0, m_0, n_0, C, F, T, N, Q$

Step 2: $n = c_0$

Step 3: Determine $x_1(q)$

Step 4: $X_1(m) = \text{FFT}[x_1(q)]$

Step 5: Convert $X_1(m)$ into $X(n\Delta_t, m\Delta_f)$

Step 6: Set $n = n+1$ and return to Step 3 until $n = c_0 + C - 1$.

Complexity = ?

IV-F Non-Uniform Δ_t

(A) 先用較大的 Δ_t

(B) 如果發現 $\left|X(n\Delta_t, m\Delta_f)\right|$ 和 $\left|X((n+1)\Delta_t, m\Delta_f)\right|$ 之間有很大的差異
則在 $n\Delta_t$, $(n+1)\Delta_t$ 之間選用較小的 sampling interval Δ_{t1}

($\Delta_\tau < \Delta_{t1} < \Delta_t$, Δ_t/Δ_{t1} 和 Δ_{t1}/Δ_τ 皆為整數)

再用 page 131 的方法算出

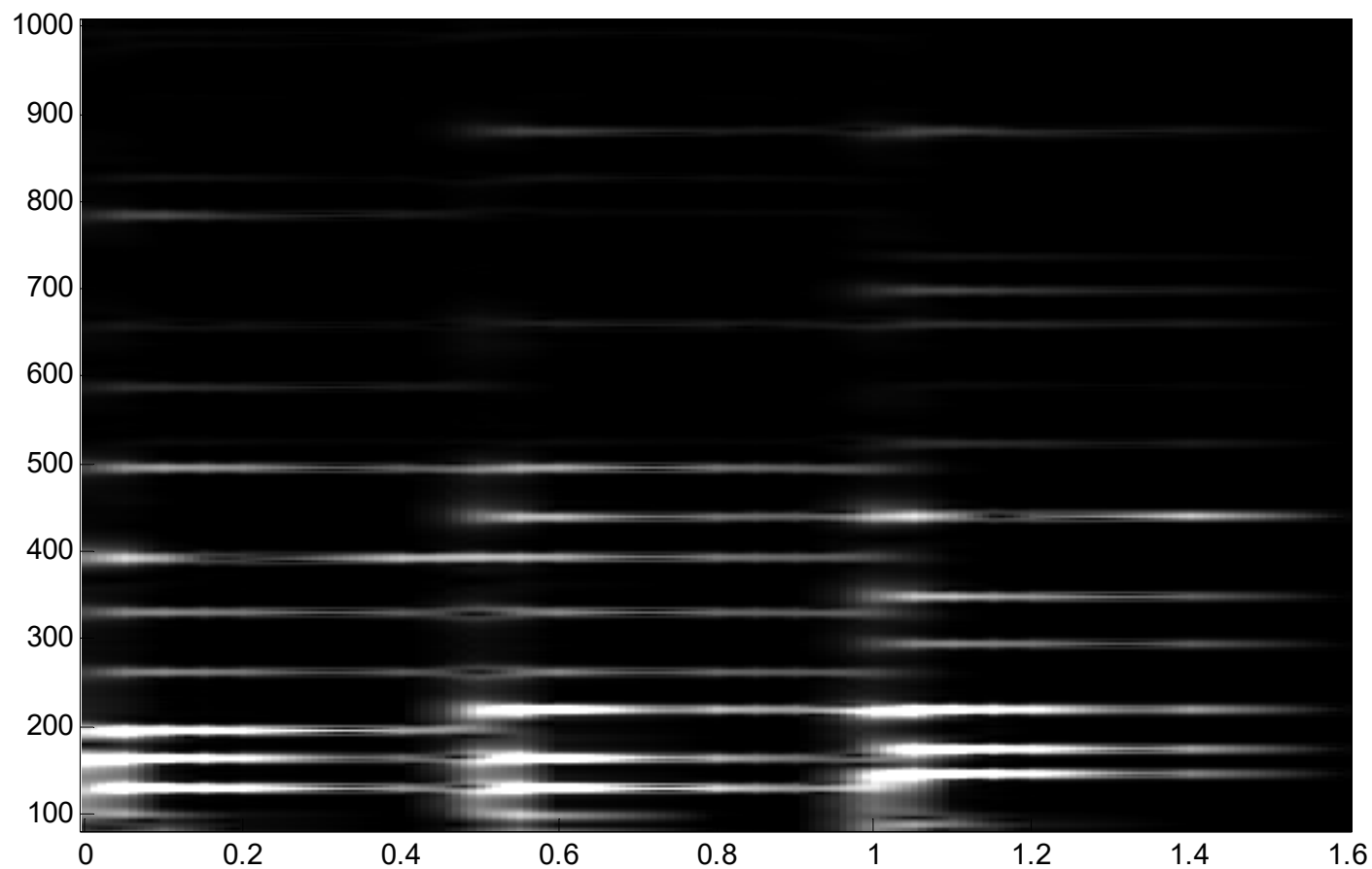
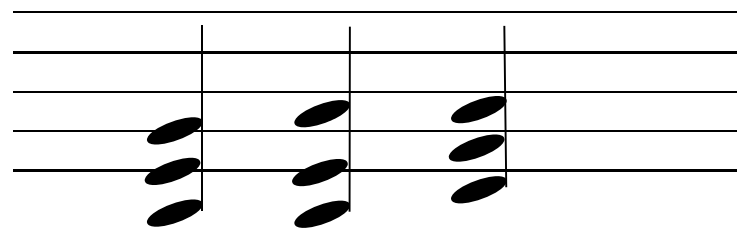
$$X(n\Delta_t + \Delta_{t1}, m\Delta_f), \quad X(n\Delta_t + 2\Delta_{t1}, m\Delta_f), \quad \dots, \quad X((n+1)\Delta_t - \Delta_{t1}, m\Delta_f)$$

(C) 以此類推，如果 $\left|X(n\Delta_t + k\Delta_{t1}, m\Delta_f)\right|$, $\left|X(n\Delta_t + (k+1)\Delta_{t1}, m\Delta_f)\right|$

的差距還是太大，則再選用更小的 sampling interval Δ_{t2}

($\Delta_\tau < \Delta_{t2} < \Delta_{t1}$, Δ_{t1}/Δ_{t2} 和 Δ_{t2}/Δ_τ 皆為整數)

Gabor transform of a music signal



$\Delta_\tau = 1/44100$ (總共有 $44100 \times 1.6077 \text{ sec} + 1 = 70902$ 點)

(A) Choose $\Delta_t = \Delta_\tau$

running time = out of memory (2008年) 15.262140 sec (2022年)

(B) Choose $\Delta_t = 0.01 = 441\Delta_\tau$ ($1.6/0.01 + 1 = 161$ points)

running time = 1.0940 sec (2008年) 0.041053 sec (2022年)

(C) Choose the non-uniform sampling points on the t -axis as

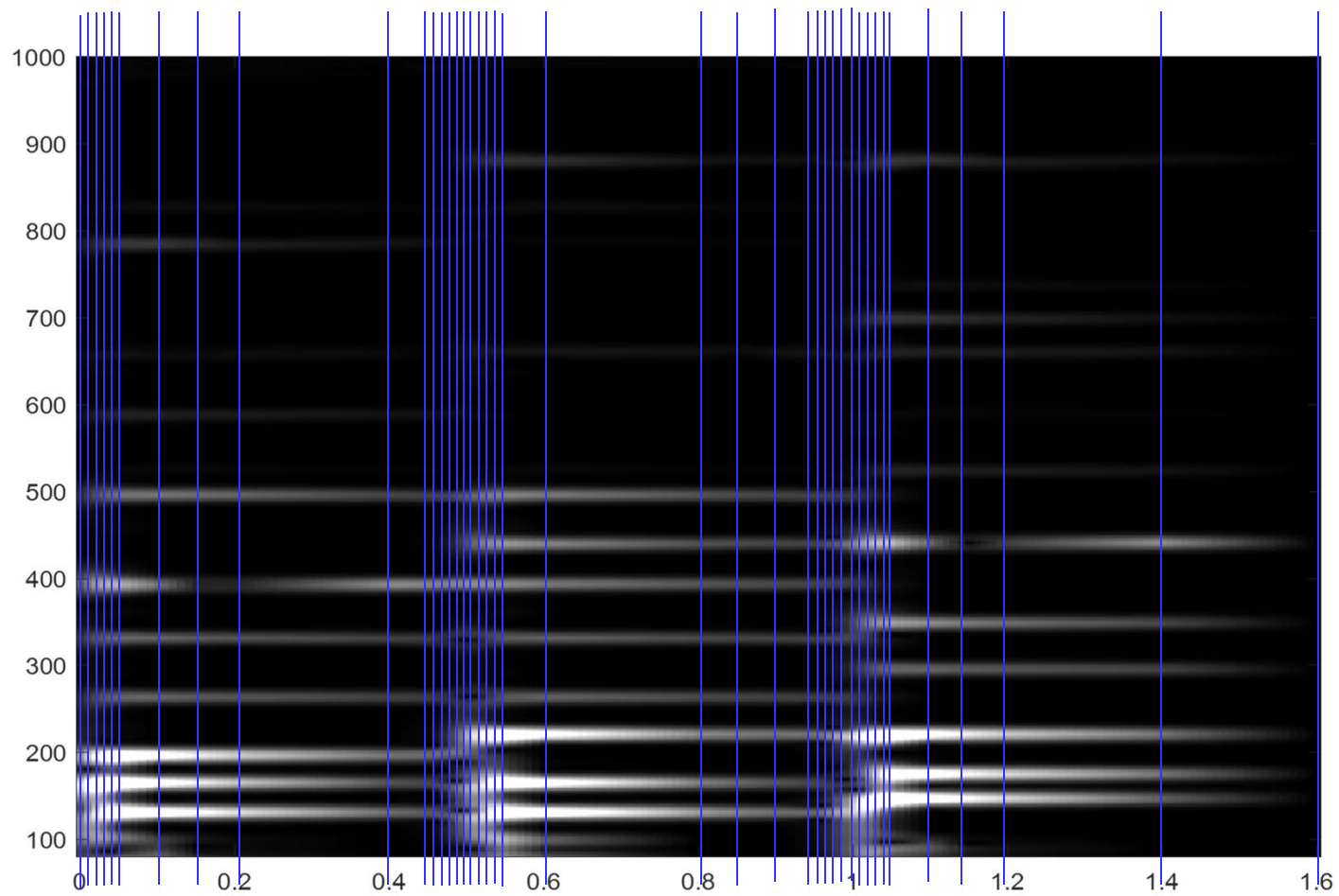
$t = 0, 0.01, 0.02, 0.03, 0.04, \underline{0.05}, \underline{0.1}, \underline{0.15}, 0.2, 0.4, \underline{0.45}, 0.46, 0.47, 0.48,$
 $0.49, \underline{0.5}, 0.51, 0.52, 0.53, 0.54, \underline{0.55}, 0.6, 0.8, \underline{0.85}, \underline{0.9}, \underline{0.95}, 0.96, 0.97,$
 $0.98, 0.99, 1, 1.01, 1.02, 1.03, 1.04, \underline{1.05}, \underline{1.1}, \underline{1.15}, 1.2, 1.4, 1.6$

(41 points)

intervals: $0.2 \rightarrow 0.05 \rightarrow 0.01$

running time = 0.2970 sec (2008年) 0.010594 sec (2022年)

with adaptive output sampling intervals



附錄七 和 Dirac Delta Function 相關的常用公式

$$(1) \int_{-\infty}^{\infty} e^{-j2\pi t f} dt = \delta(f)$$

$$(2) \delta(t) = |a| \delta(at) \quad (\text{scaling property})$$

$$(3) \int_{-\infty}^{\infty} e^{-j2\pi t g(f)} dt = \delta(g(f)) = \sum_n |g'(f_n)|^{-1} \delta(f - f_n)$$

where f_n are the zeros of $g(f)$

$$(4) \int_{-\infty}^{\infty} \delta(t - t_0) y(t, \dots) dt = y(t_0, \dots) \quad (\text{sifting property I})$$

$$(5) \delta(t - t_0) y(t, \dots) = \delta(t - t_0) y(t_0, \dots) \quad (\text{sifting property II})$$